

CSCI1680

Network Layer: IP & Forwarding

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Administivia

- Snowcast: due Wednesday (9/27)
- IP project: out Thursday, fill out group preference form by Thursday 11:59pm
- ↳ GEARUP THURS 5-7 PM + ZOOM + REL.
- **HW1: out later today, due next Thurs**
 - Some practice for IP!

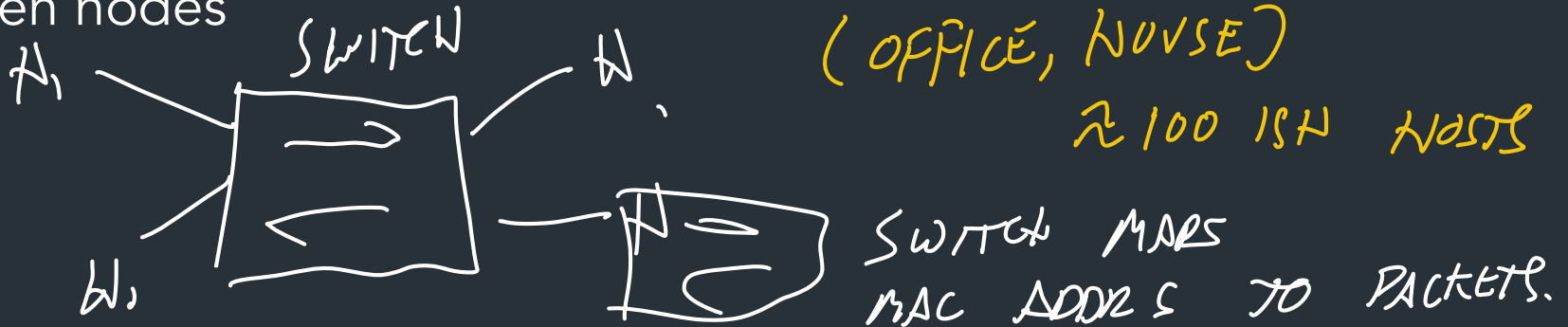
Today

Start of network layer

- Network layer: Internet Protocol (IP) (v4)
- Mechanics of IP forwarding
- Intro to IP project

Last time: link-layer

- Hosts connect to network via interfaces
- Every interface has a link-layer address
Ethernet/Wifi: MAC address (0c:45:22:c1:be:03)
- Switches, Wifi APs: in-network devices that forward packets between nodes



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Mental model for the link layer

=> How to connect hosts on a "small" network

=> Given link-layer address, know how to reach host on your network

Layers, Services, Protocols

Application

Service: user-facing application.
Application-defined messages

Transport

Service: multiplexing applications
Reliable byte stream to other node (TCP),
Unreliable datagram (UDP)

L3

Network

Service: move packets to any other node in the network
Internet Protocol (IP)

L2

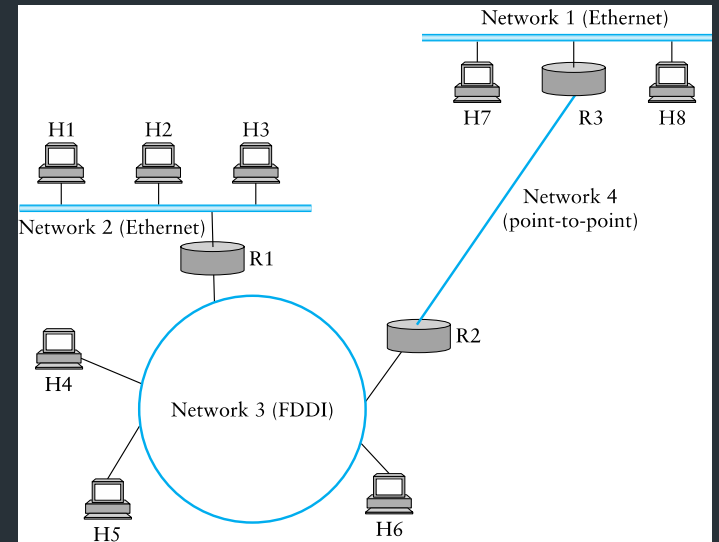
Link

Service: move frames to other node across link.
May add reliability, medium access control

Physical

Service: move bits to other node across link

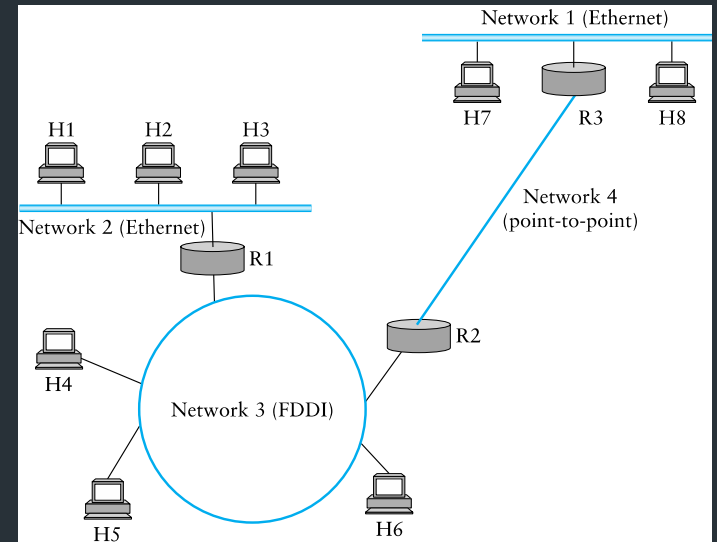
Internet Protocol (IP) Goals



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How to connect everyone?

- Glue lower-level networks together
- A network of networks!
- Router: device that forwards packets between networks



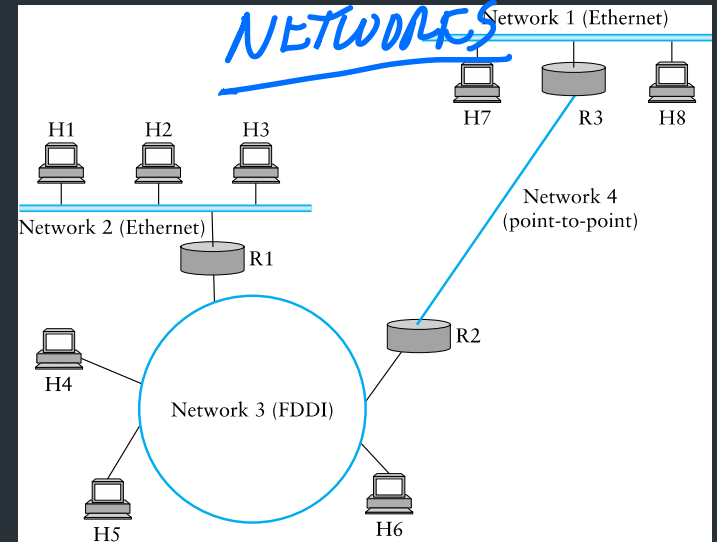
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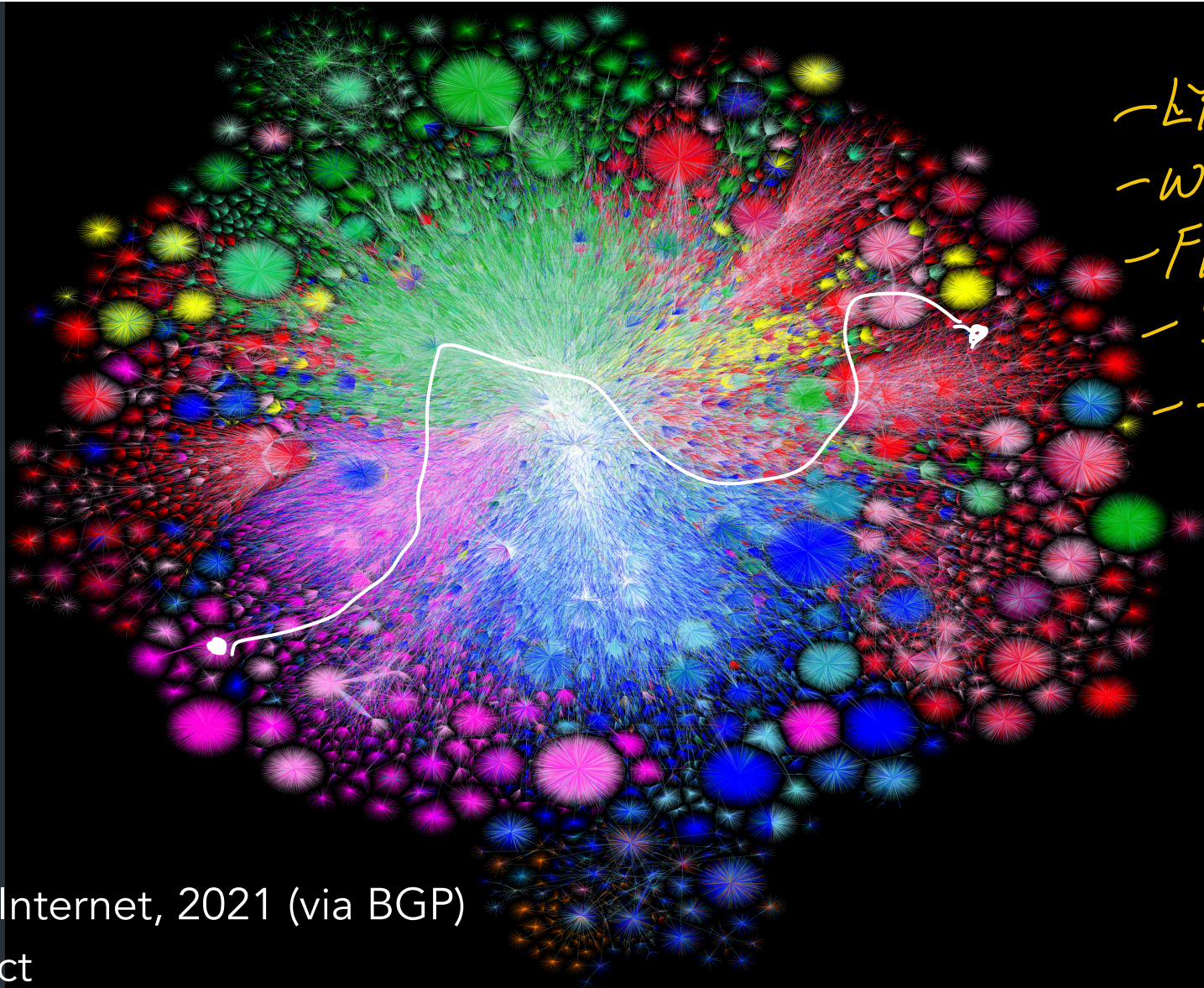
How to connect everyone?

- Glue lower-level networks together
- A network of networks!
- Router: device that forwards packets between networks

=> Doesn't that sound like switching?

*LOOKS AT DIFFERENT
INFO. MOVES BETWEEN
NETWORKS*





- ETN
- WIFI
- FIBER
- SATELLITE
- - - -

Color Chart

North America (ARIN)	
Europe (RIPE)	
Asia Pacific (APNIC)	
Latin America (LANIC)	
Africa (AFRINIC)	
Backbone	
US Military	

Map of the Internet, 2021 (via BGP)
 OPTE project

New Challenges

- Networks are heterogeneous (eg. Wifi vs. Ethernet)
 - Different frame formats
 - Some are more reliable than others
 - Different packet sizes/bandwidths

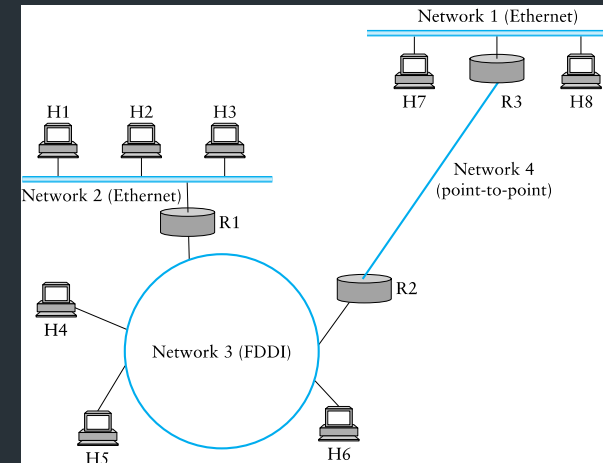
CAN'T MAKE TOO MANY ASSUMPTIONS ABOUT L2

- Scaling: link-layer strategies don't work!

SWITCH: TABLE

MAC ADDR \Rightarrow PORT

NEED DIFF STRATEGIES!



What came before the Internet?

The (landline) telephone network



(Plain Old Telephone Service (POTS))

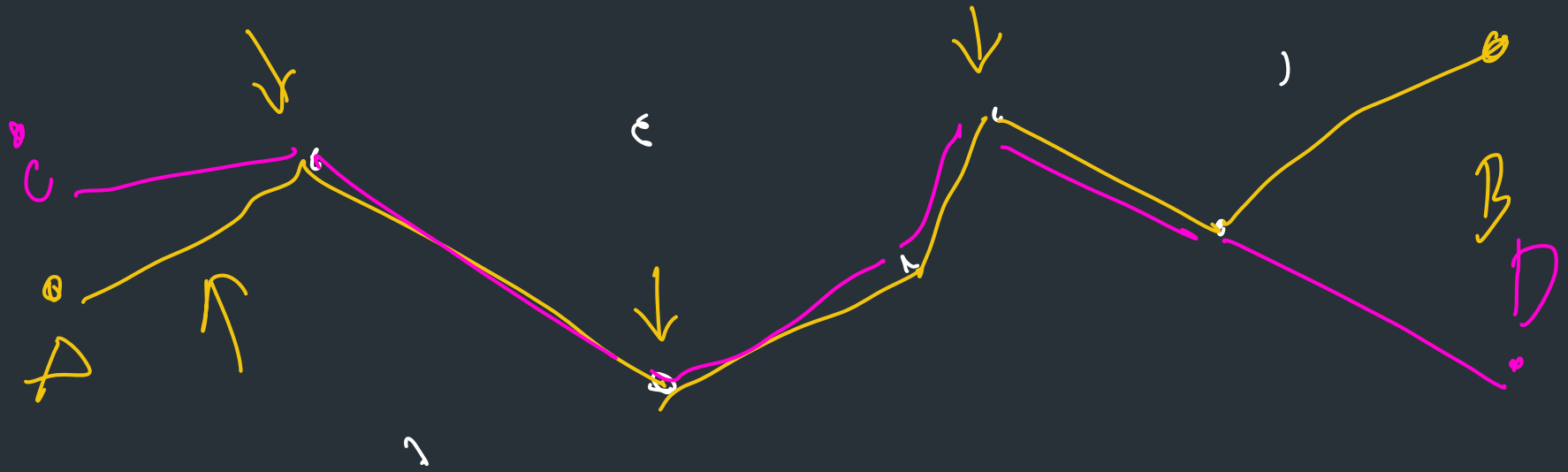
At the time the internet was forming (~1960s), landline phones were a big network that connected lots of end-users (160M worldwide)

Why should we care?

⇒ Provides a useful comparison

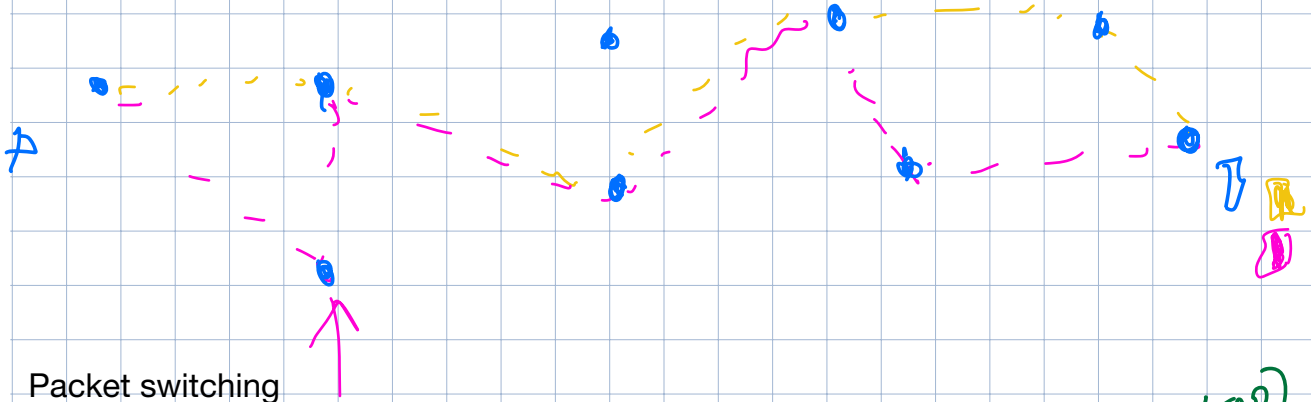
⇒ Understand some of the design goals of the time

EARLY PHONE NETWORKS: CIRCUIT SWITCHING



- SETUP CONNECTION BEFORE YOU TALK
- NODES NEED TO REMEMBER CONNECTIONS

PACKET SWITCHING: ROUTERS MOVE SMALL CHUNKS OF DATA ("PACKETS")



Packet switching

- Routers just need to pick destination for next packet
- No initial setup
- Messages must get divided up into small units (and reassembled at other end)
- Packets may take any path (can be different for each packet)

(TCP)

Early telephone networks

Used circuit switching: set up whole path for call beforehand

- Does it scale?

WORKED FOR TELEPHONES UNTIL 1980s,
BUT INTERNET HAS VERY DIFFERENT DESIGN GOALS.

A Bit of History

Early Packet switched networks: Arpanet's IMPs

- Late 1960's => RFC 1, 1969!
 - Reliable network with many features we know today
-
- Build apps with Network Control Program (NCP)
 - Built on reliable IMPs
 - Used by programs like telnet, mail, file transfer



A Bit of History

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What about when network isn't reliable?

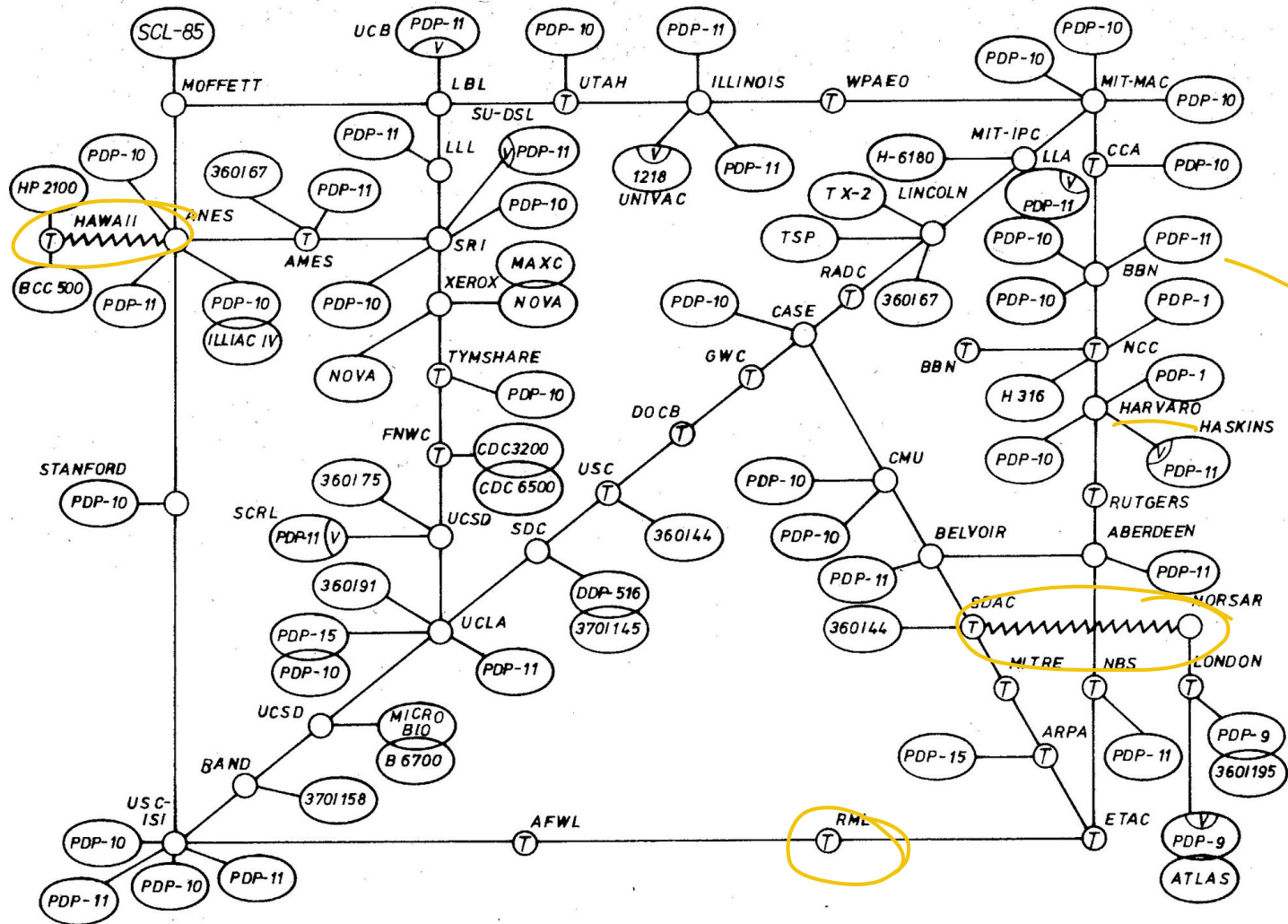


CS/360



PDP 10
(CONNECTING
SINGLE
COMPUTERS!!)

THE ARPA NETWORK
DEC 1969



SINGLE
COMPUTER

Abb. 4 ARPA Network, topologische Karte. Stand Juni 1974.

How to make such a protocol?

- How to deal with heterogeneous networks?
- How to find hosts?
- Should messages be reliable or unreliable?
- What to do when a device joins/leaves?
- ...

NEED TO WORK
W/ "LEAST COMMON
DENOMINATOR"
OF FUNCTIONALITY"

Big concerns

⇒ Not every application needs all features

⇒ Can't assume much functionality from (heterogeneous) link layer

1974: TCP/IP Introduced

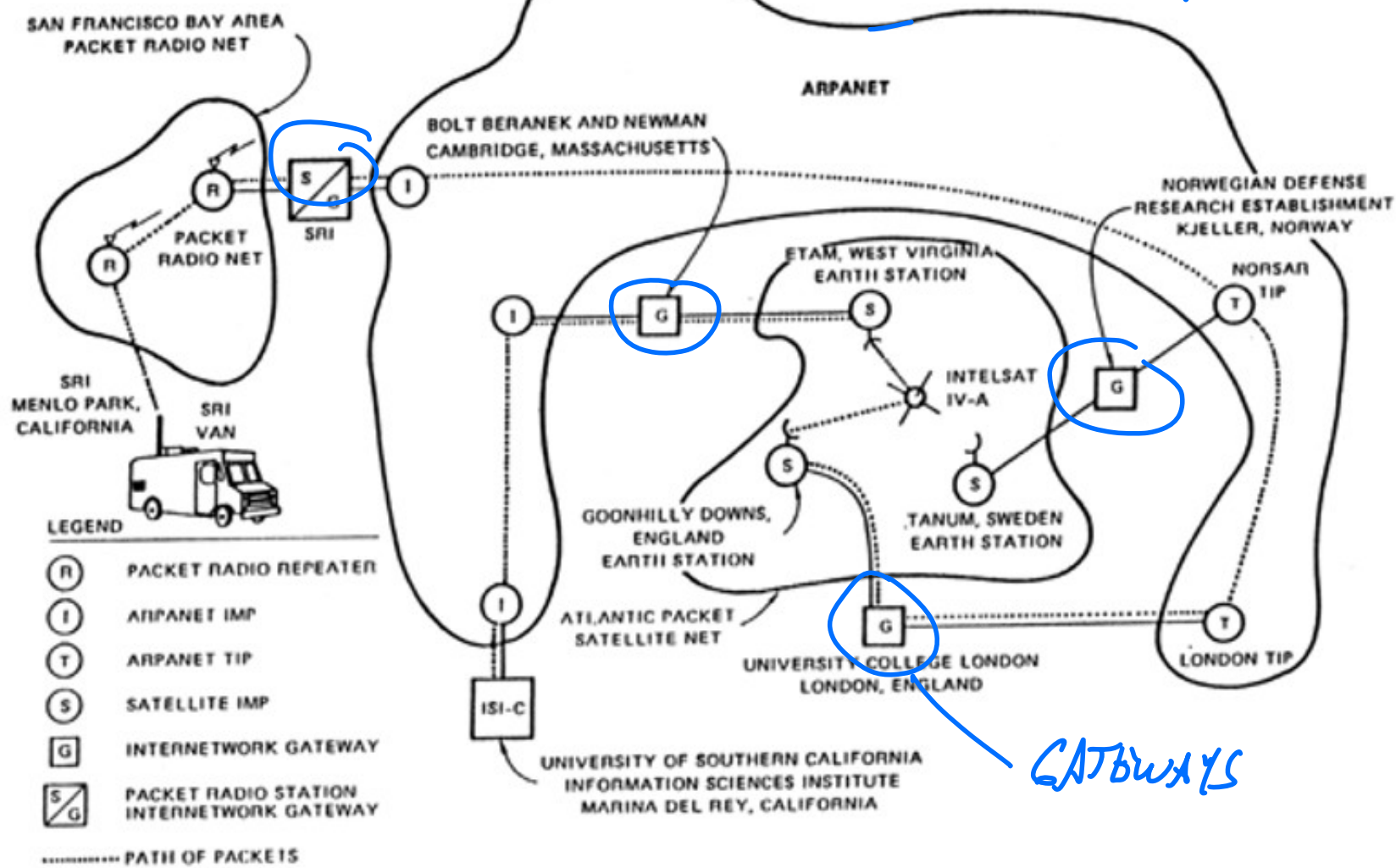
- Vint Cerf, Robert Kahn build protocol to replace NCP
- Initial design: single protocol providing a reliable pipe

1974: TCP/IP Introduced

- Vint Cerf, Robert Kahn build protocol to replace NCP
 - Initial design: single protocol providing a reliable pipe

 - Eventually, separated into different protocols we know today
 - IP: basic datagram service among hosts
 - TCP: reliable transport
 - UDP: unreliable multiplexed datagram service
- "PACKET" OR "MESSAGE"*

-DIFF NET TYPES



IP's Decisions

- Connectionless, packet-switched network

≡ NO CONNECTION SETUP.

- "Best-effort" service: other layers add reliability if you need

it

*NO GUARANTEES ON RELIABILITY w/
IP ALONE.*

How to reach hosts?

- Common message format: IP header
- Every host identified by an IP address



IP's Decisions

- Connectionless, packet-switched network
 - => Routers are "simple" => no connection state
- "Best-effort" service: other layers add reliability if you need it
 - => Packets might be dropped, reordered, delayed, ...

How to reach hosts?

- Common message format: IP header
- Every host identified by an IP address

An excellent read

David D. Clark, "The design Philosophy of the DARPA Internet Protocols",
1988

- Primary goal: multiplexed utilization of existing interconnected networks
- Other goals:
 - Communication continues despite loss of networks or gateways
 - Support a variety of communication services
 - Accommodate a variety of networks
 - Permit distributed management of its resources
 - Be cost effective
 - Low effort for host attachment
 - Resources must be accountable

The Internet Protocol

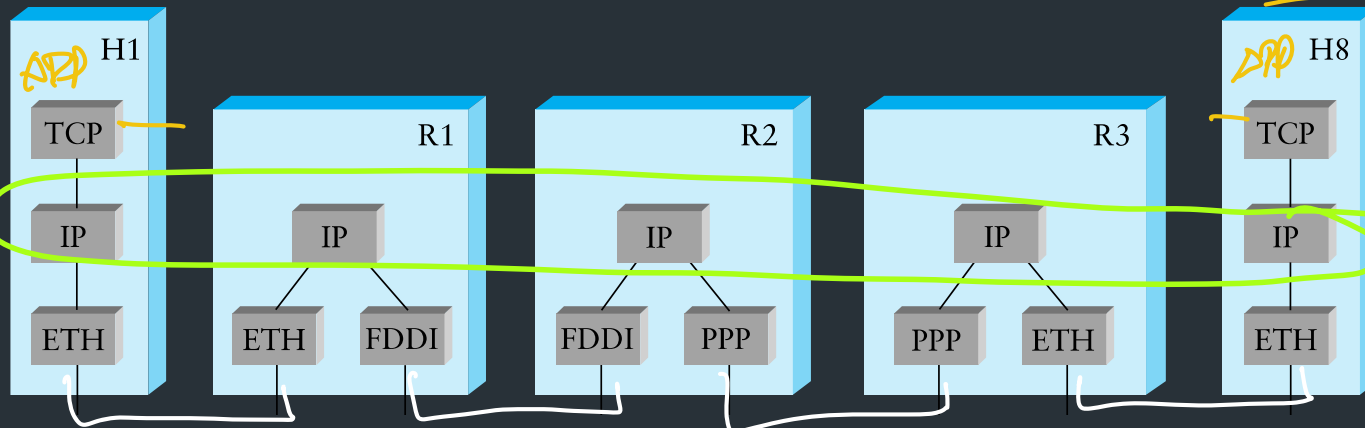
IP runs on all hosts and routers

- Addressing: how we name nodes in an IP network
- Provides forwarding: how routers move packets based on the destination address
- (later) Routing: how routers build forwarding rules

DATA PLANE

CONTROL PLANE

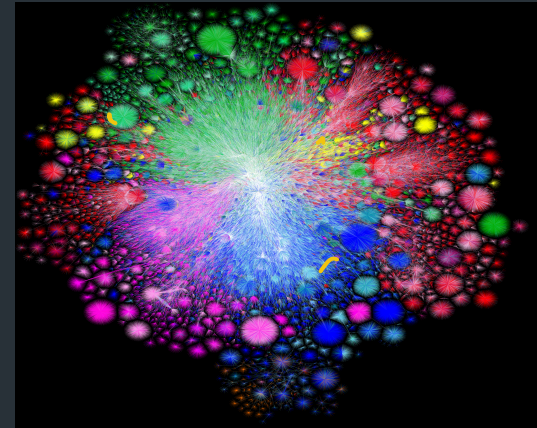
USED BY EVERY NODE IN THE NETWORK!



IP Addressing

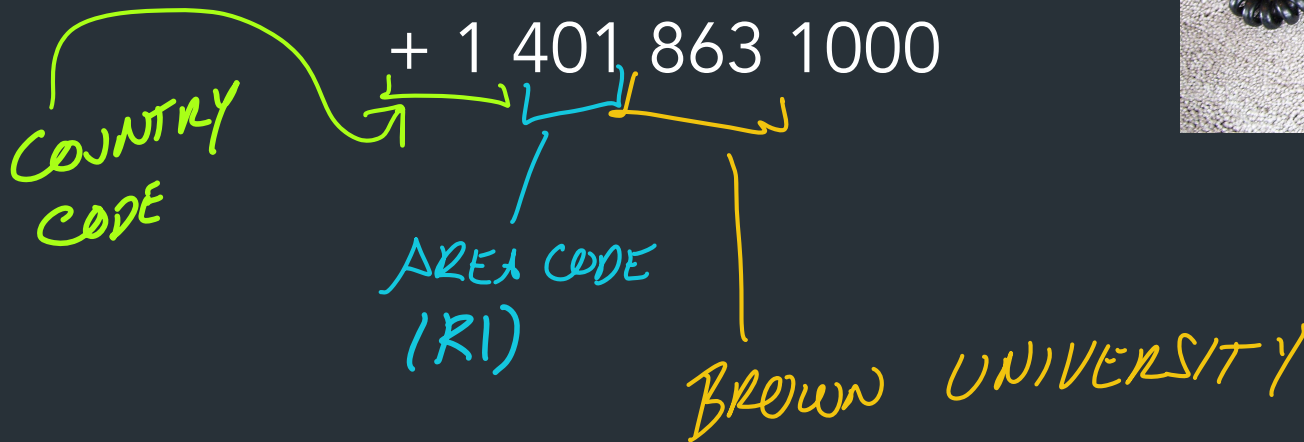
Overview

- Unique number to identify “all” hosts on the Internet
- A number with structure => the number tells the network where the host is



Analogy: back to phones

Telephone numbers have a structure to them



Analogy: back to phones

Telephone numbers have a structure to them



+ 1 401 863 1000



RI

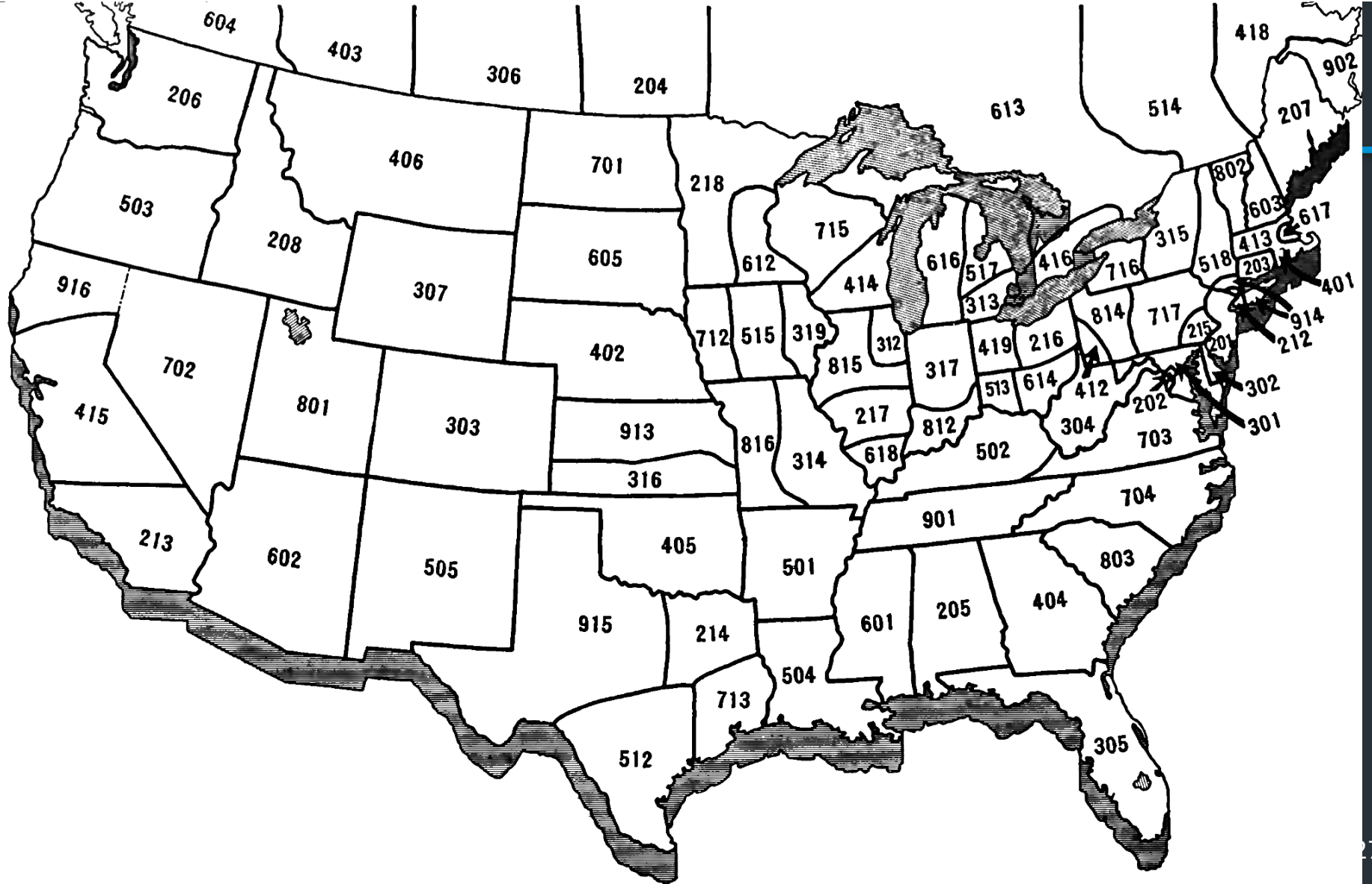
+1 212 555 4253

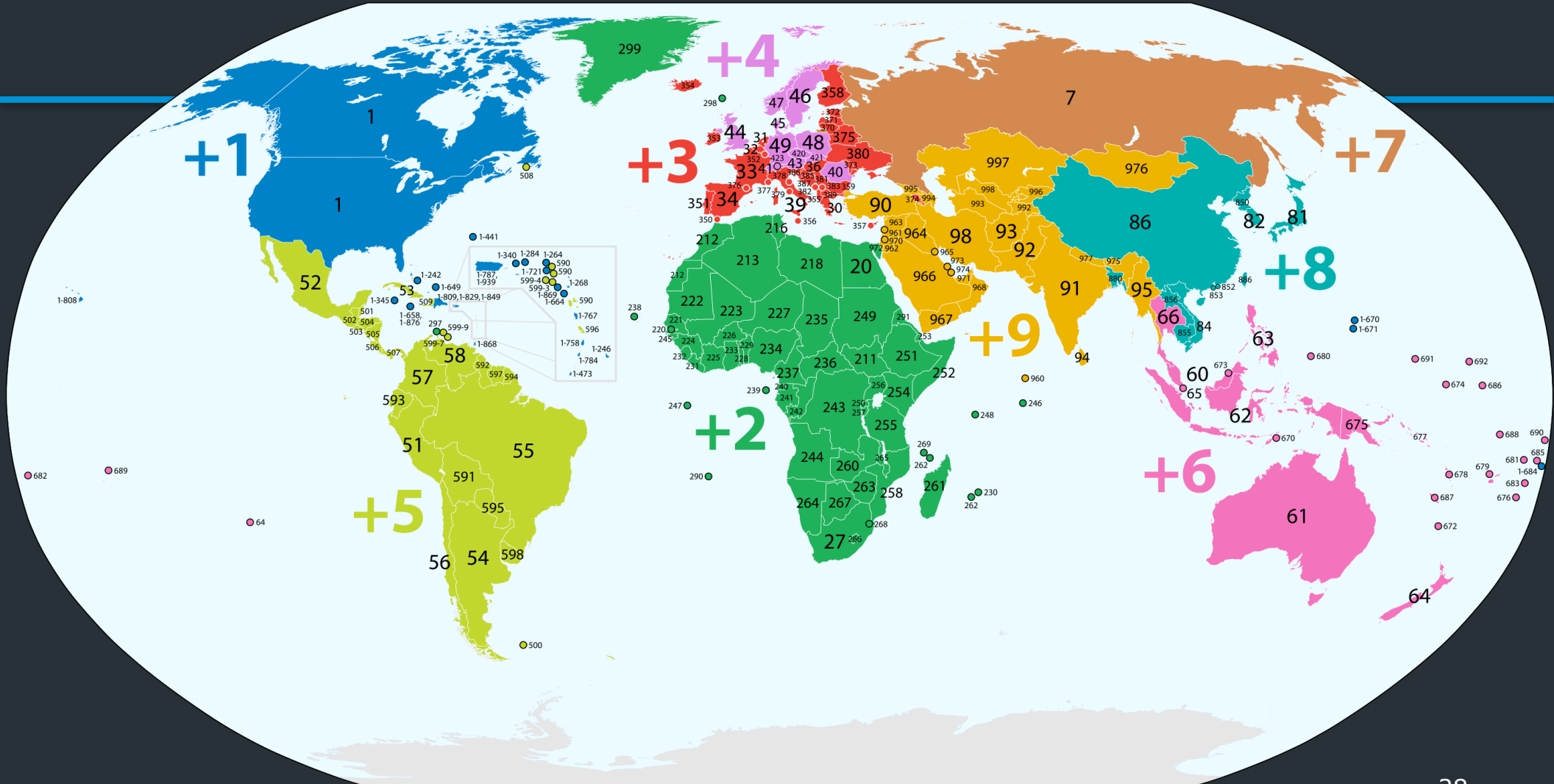
MANHATTAN



⇒ AT ONE TIME,
USED TO ROUTE CALLS
TO A GEOGRAPHIC
LOCATION.

Part of the number tells ~~where~~ you are!
(or at least it did before cell phones)





IP Addressing



IP Version 4: Each address is a 32-bit number:

128.148.16.7

128.148.16.7

10000000 10010100 00010000 00000111

32 bits => 2^{32} possible addresses...
problem?

Notation

- Write each byte ("octet") as a decimal number
- This is called "dotted decimal" or "dotted quad" notation

IP Addressing

An IP address identifies...

- Who a host is: A unique number
- Where it is on the Internet
- Networks are allocated ranges of IPs by global authority (ICANN)
 - Further subdivided by regions, ISPs, ...
- Some IPs have special uses (eg. 127.0.0.1)

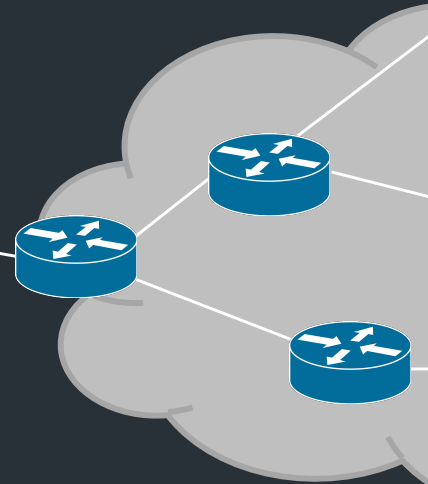
→ HOST #1247

↳ IDENTIFIES THE NETWORK

(EG. Brown, Amazon, Google, ...)

128.148.16.7

↳ LOCAL HOST



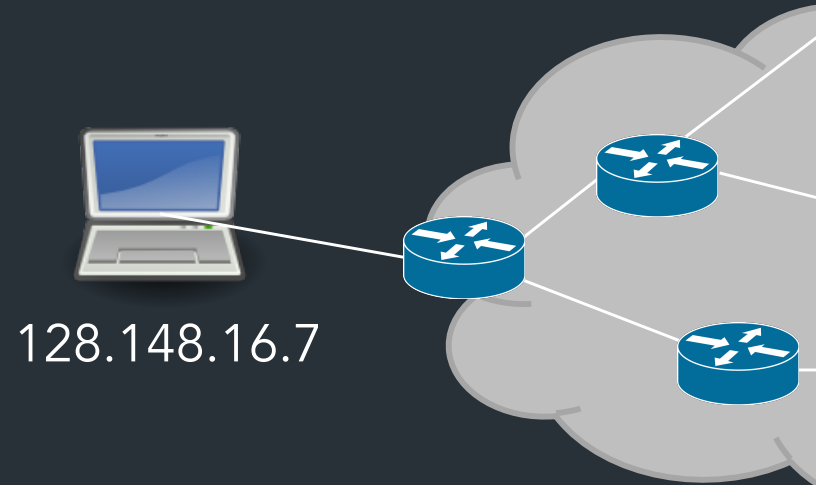
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eg. Brown owns 128.148.xxx.xxx, 138.16.xxx.xxx



IP Addressing

Brown owns
128.148.0.0/16

Brown owns the range:

128.148.xxx.xxx

16 BITS

$2^{16} \approx 65K$ HOSTS

10000000 10010100 xxxxxxxxxx xxxxxxxxxx

16 BITS

Network part

Identifies Brown (to the Internet)

INDIVIDUAL HOSTS

Host part

Denotes individual hosts within the Brown Network

⇒ USED BY INTERNET TO
SEND PACKETS TO THIS NETWORK.

Formal way to write this:

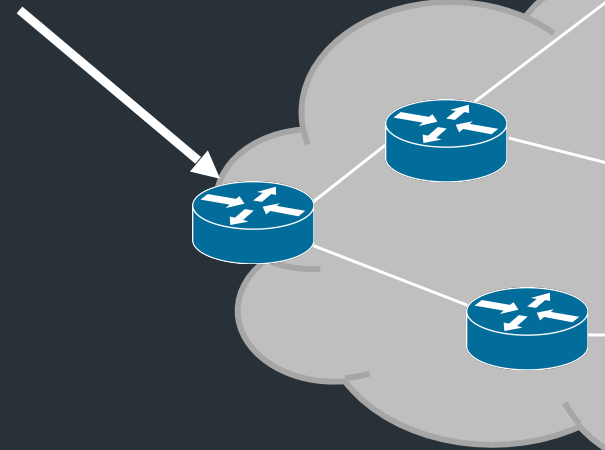
128.148.0.0/16

IP Addressing

A network can designate IP addresses for its own hosts within its address range

For 128.148.xxx.xxx:

Brown



IP Addressing

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For 128.148.xxx.xxx:

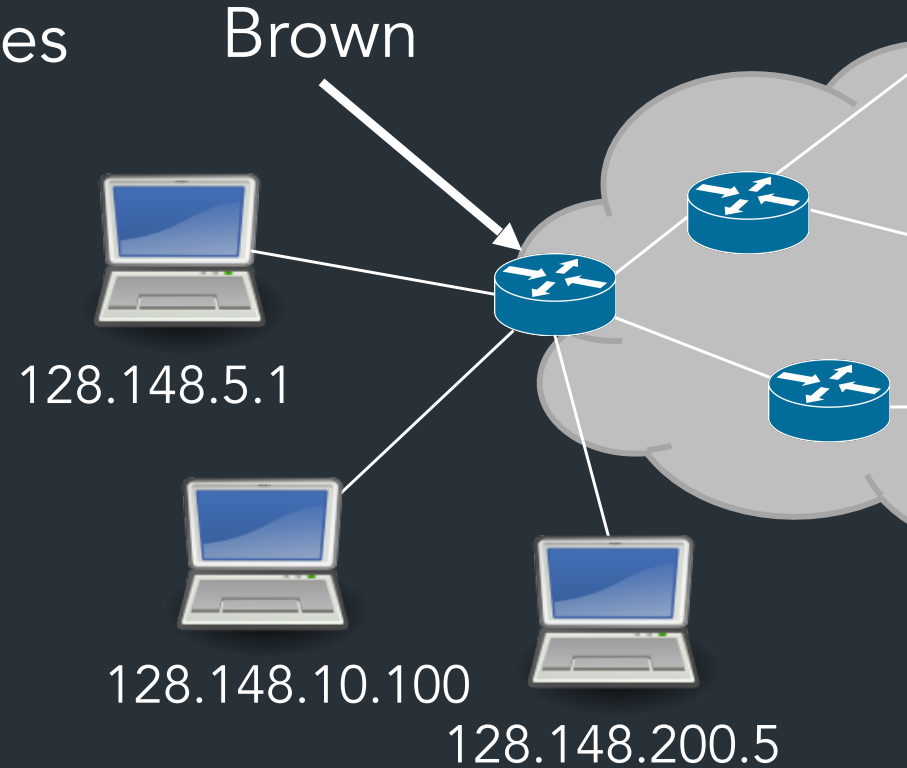
10000000 10010100 xxxxxxxx xxxxxxxx

Brown uses the the prefix 128.148.0.0/16

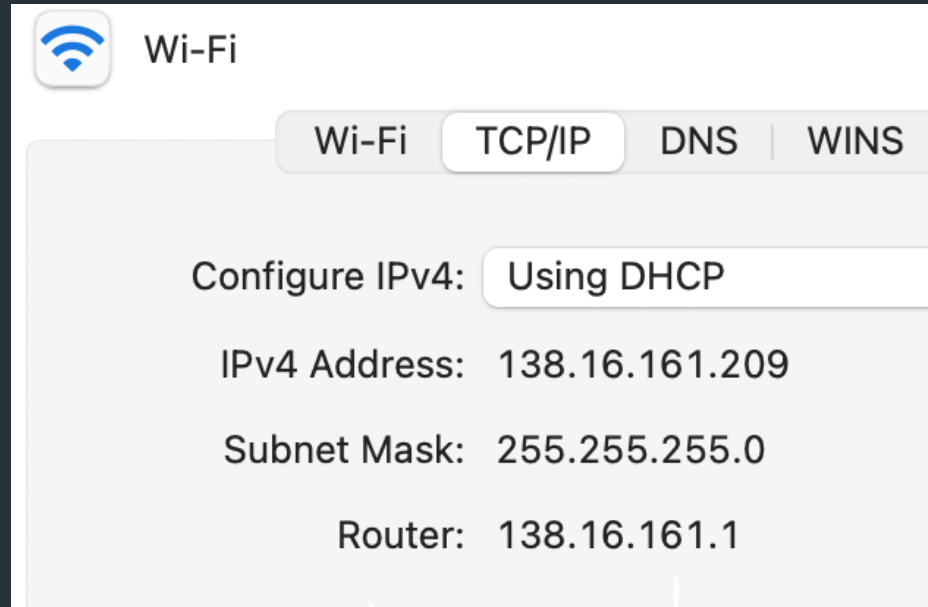
Some other ways to write this:

128.148/16

128.148.0.0 + subnet mask 255.255.0.0



A typical configuration



IP ADDRESS

138
1000 1010

16
0001 0000

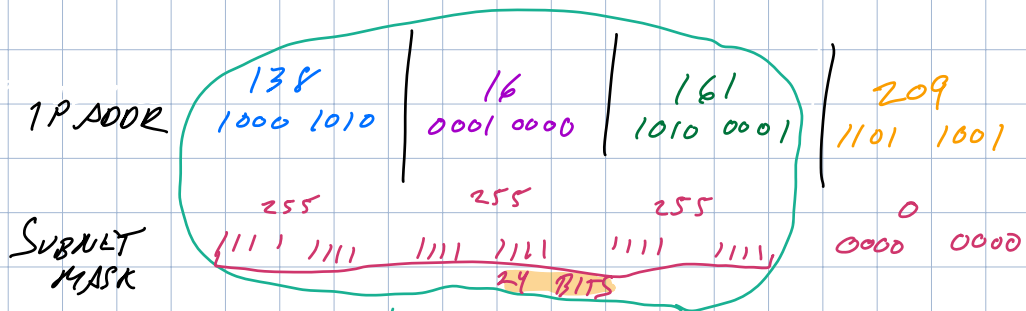
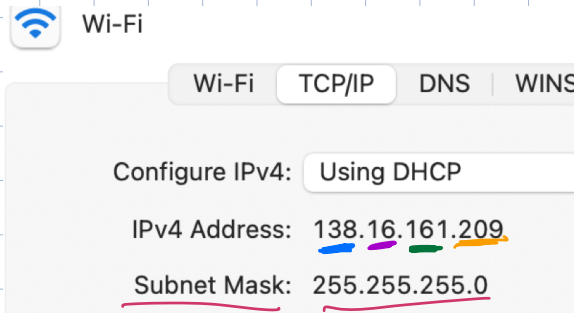
161
1010 0001

209
1101 1001

SUBNET MASK

138.16.161.0/24

UNPACKING IP SETTINGS: EXAMPLE



→ BITS WITH ONES IN SUBNET MASK IDENTIFY THE NETWORK PART.
SO THIS IS NETWORK 12

138.16.161.0/24

16, THE SUBNET MASK IS 24 BITS LONG.

Common prefixes

1.2.0.0/16

00000001 00000010 xxxxxxxxxxxx xxxxxxxxxxxx

Host PART: 16 BITS

$2^{16} \approx 65K$ HOSTS

8.0.0.0/8

00001000 xxxxxxxxxxxx xxxxxxxxxxxx xxxxxxxxxxxx

Host PART = $2^{32-8} = 2^{24} \approx 16M$ HOSTS

123.10.1.0/24

01111011 00001010 00000001 xxxxxxxxxx

$2^8 = 256$ HOSTS

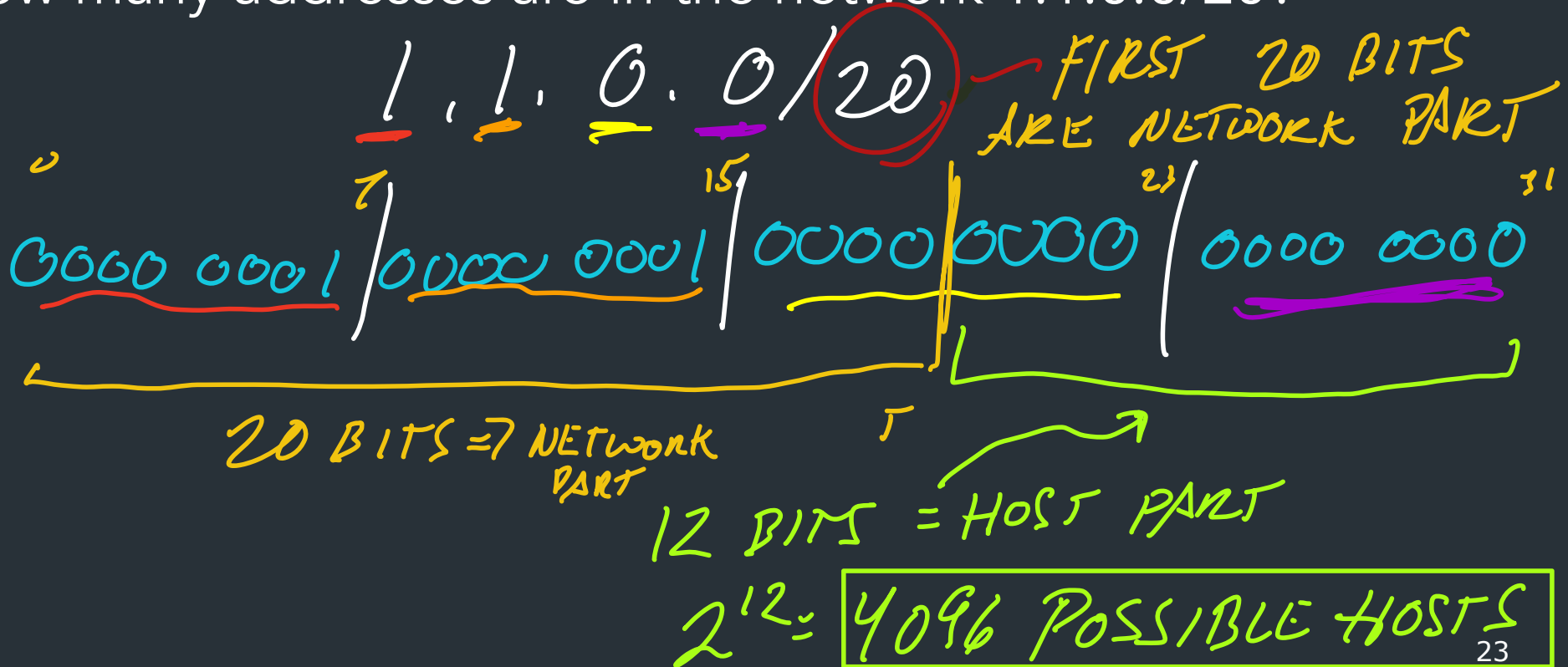
201.112.10.200/30

11001001 01110000 00001010 110010xx

$2^2 = 4$ HOSTS

Example

How many addresses are in the network 1.1.0.0/20?



IP PREFIX:

A RANGE OF ADDRESSES
FOR A NETWORK.

How do we move
packets between networks?

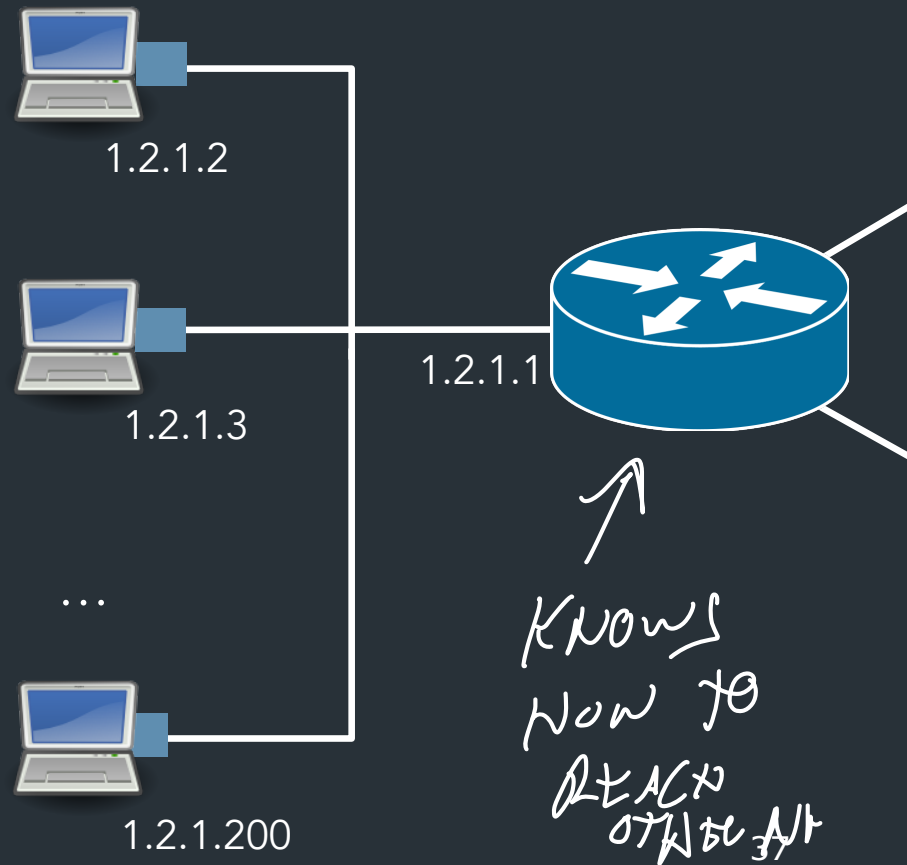
How IP forwarding works

Assume:

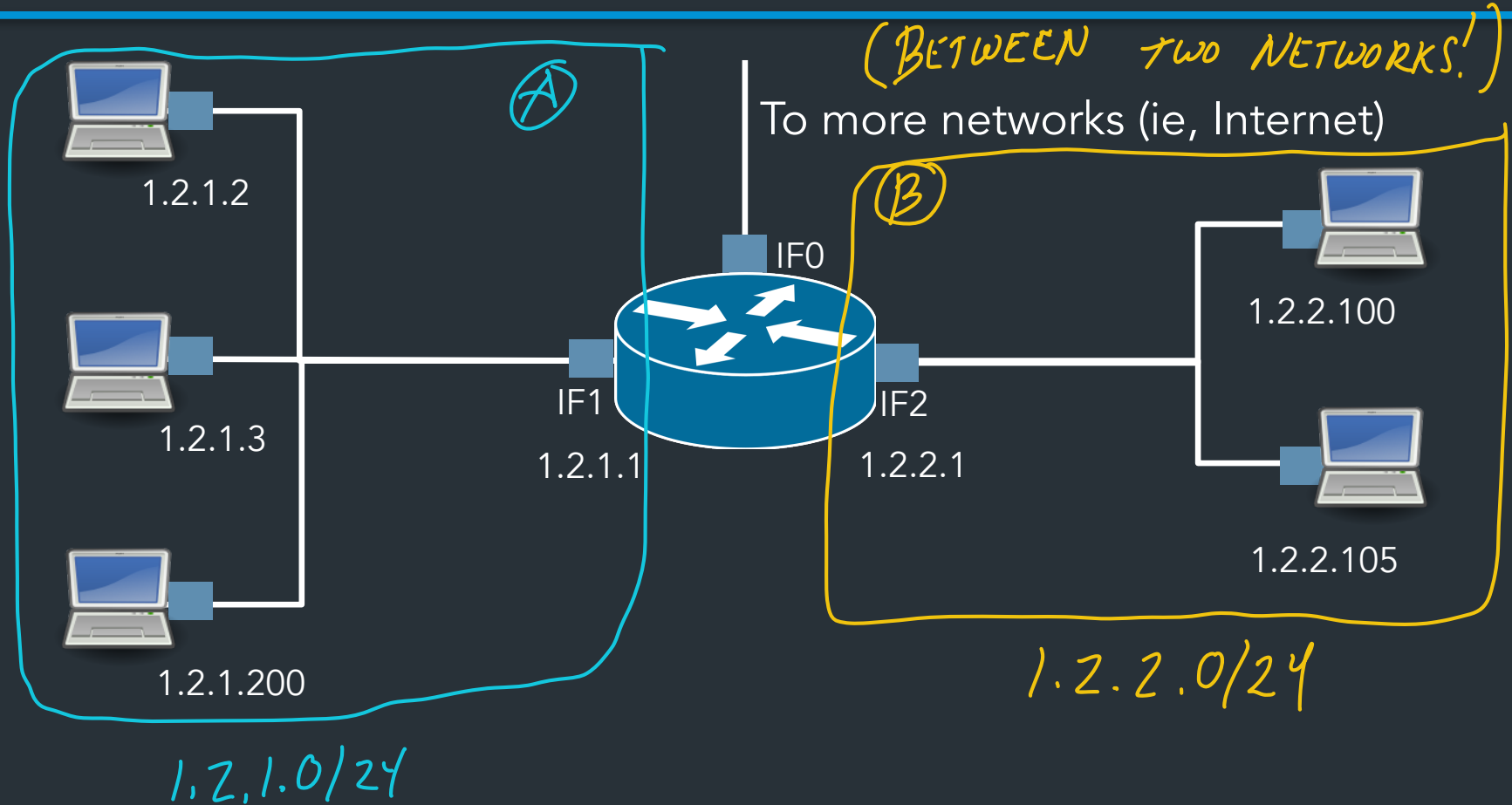
- Communicating on same network is easy—this is the link-layer's job!
- Can map IP addresses to MAC addresses (more on this later)

How to reach an address outside this

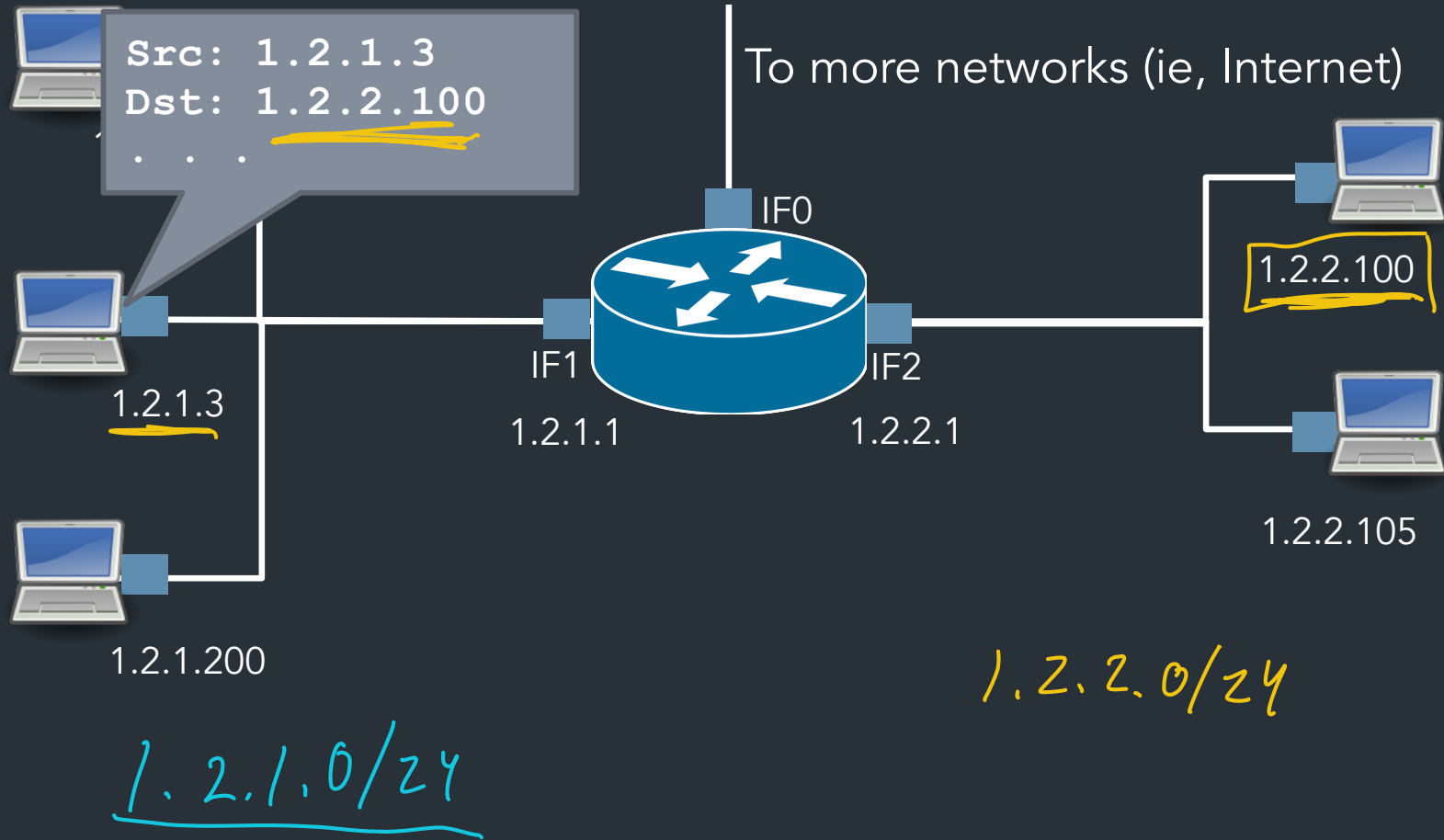
Send packets to a router,
which forwards IP packets
to other networks



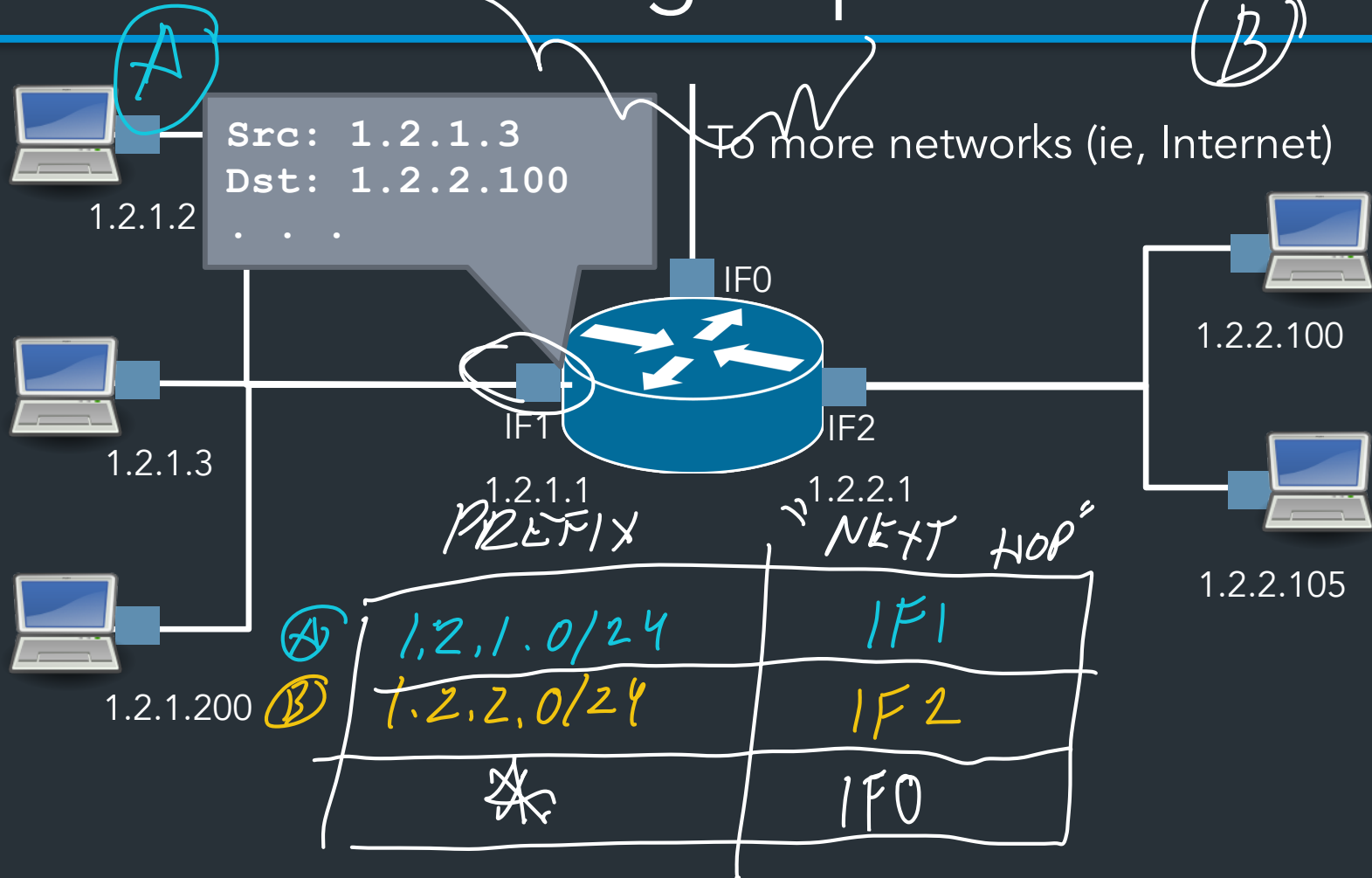
Forwarding IP packets



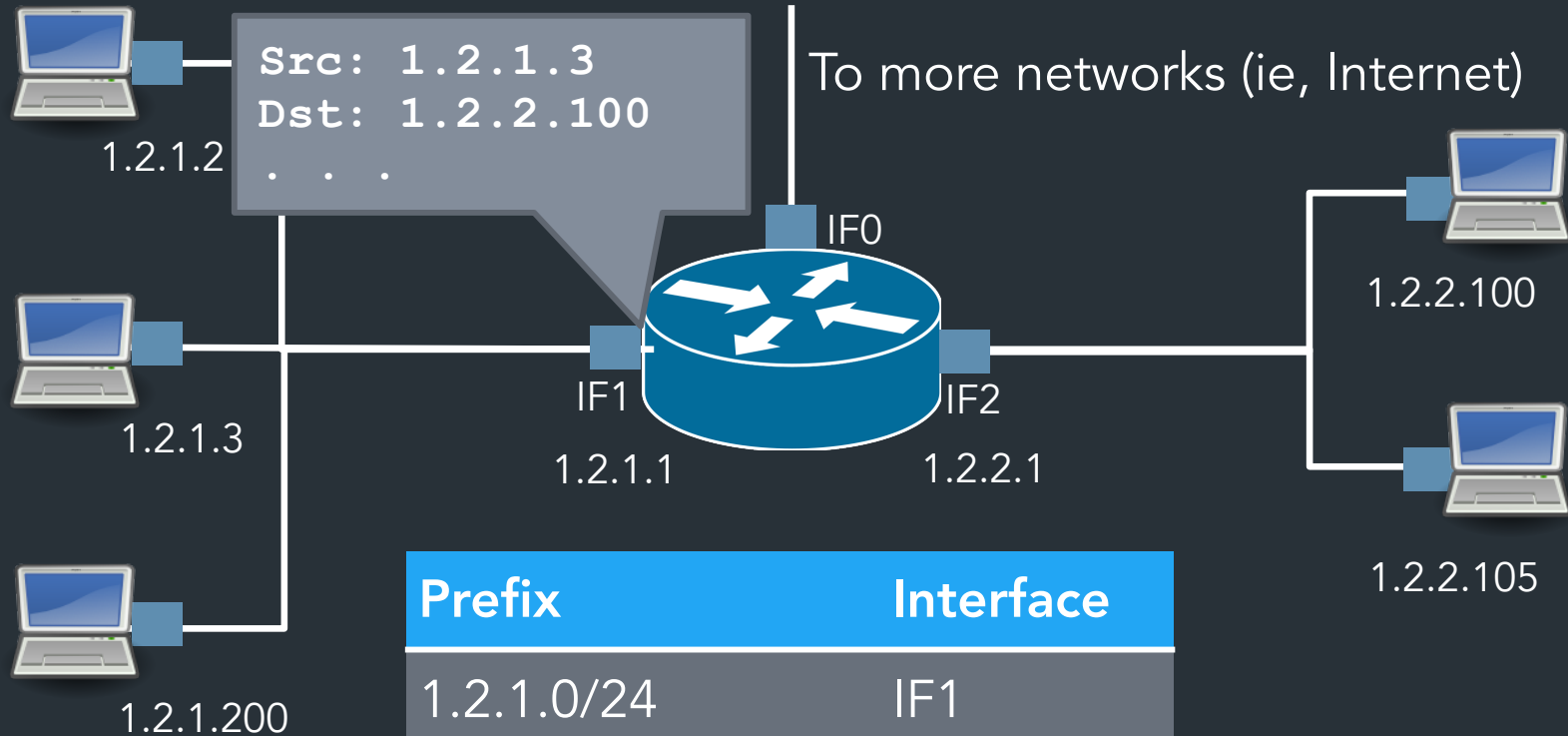
Forwarding IP packets



Forwarding IP packets



Forwarding IP packets



Prefix	Interface
1.2.1.0/24	IF1
1.2.2.0/24	IF2
<everything else>	(IF0)

A routing table

```
R6#sh ip ro
Gateway of last resort is 108.34.215.1 to network 0.0.0.0

S*   0.0.0.0/0 [1/0] via 108.34.215.1
      10.0.0.0/8 is variably subnetted, 7 subnets, 3 masks
C     10.1.0.0/24 is directly connected, wlan-ap0
L     10.1.0.2/32 is directly connected, wlan-ap0
O IA  10.1.44.1/32 [110/1001] via 10.20.30.33, 3w4d, Tunnel0
C     10.1.48.0/24 is directly connected, Loopback0
L     10.1.48.1/32 is directly connected, Loopback0
C     10.20.30.32/31 is directly connected, Tunnel0
L     10.20.30.32/32 is directly connected, Tunnel0
      108.0.0.0/8 is variably subnetted, 2 subnets, 2 masks
C     108.34.215.0/24 is directly connected, GigabitEthernet0/0
L     108.34.215.208/32 is directly connected, GigabitEthernet0/0
      172.16.0.0/16 is variably subnetted, 2 subnets, 2 masks
C     172.16.98.0/24 is directly connected, Vlan98
L     172.16.98.1/32 is directly connected, Vlan98
      172.17.0.0/16 is variably subnetted, 6 subnets, 3 masks
O IA  172.17.44.0/24 [110/1001] via 10.20.30.33, 3w4d, Tunnel0
C     172.17.48.0/24 is directly connected, Vlan20
```

A routing table

```
R6#sh ip ro
```

```
Codes: L - local, C - connected, S - static, R - RIP, M - mobile, B - BGP  
D - EIGRP, EX - EIGRP external, O - OSPF, IA - OSPF inter area  
N1 - OSPF NSSA external type 1, N2 - OSPF NSSA external type 2  
E1 - OSPF external type 1, E2 - OSPF external type 2  
i - IS-IS, su - IS-IS summary, L1 - IS-IS level-1, L2 - IS-IS level-2  
ia - IS-IS inter area, * - candidate default, U - per-user static route  
o - ODR, P - periodic downloaded static route, H - NHRP, l - LISP  
+ - replicated route, % - next hop override
```

```
Gateway of last resort is 108.34.215.1 to network 0.0.0.0
```

```
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C 10.1.48.0/24 is directly connected, Loopback0  
L 10.1.48.1/32 is directly connected, Loopback0
```

A large table

```
rviews@route-server.ip.att.net>show route table inet.0 active-path
```

```
inet.0: 866991 destinations, 13870153 routes (866991 active, 0 holddown, 0 hidden)  
+ = Active Route, - = Last Active, * = Both
```

```
0.0.0.0/0          *[Static/5] 5w0d 19:43:09  
                  > to 12.0.1.1 via em0.0  
1.0.0.0/24        *[BGP/170] 1d 10:24:47, localpref 100, from 12.122.83.238  
                  AS path: 7018 3356 13335 I, validation-state: valid  
                  > to 12.0.1.1 via em0.0  
1.0.4.0/22        *[BGP/170] 1d 10:24:47, localpref 100, from 12.122.83.238  
                  AS path: 7018 3356 4826 38803 I, validation-state: valid  
                  > to 12.0.1.1 via em0.0  
1.0.4.0/24        *[BGP/170] 1d 10:24:47, localpref 100, from 12.122.83.238  
                  AS path: 7018 3356 4826 38803 I, validation-state: valid  
                  > to 12.0.1.1 via em0.0  
1.0.5.0/24        *[BGP/170] 1d 10:24:47, localpref 100, from 12.122.83.238  
                  AS path: 7018 3356 4826 38803 I, validation-state: valid  
                  > to 12.0.1.1 via em0.0  
1 0 6 0/24        *[BGP/170] 1d 10:24:47 localpref 100 from 12 122 83 238
```