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Based partly on lecture notes by Rodrigo Fonseca, Scott Shenker and John Jannotti

Administrivia

- You should have done your milestone II meeting
- You have one week from today to finish TCP. Do not wait until the end.
- Final project info: Thursday

HTTP: Hypertext Transfer Protocol



"Application protocol for distributed, collaborative hypermedia information systems"



"Application protocol for distributed, collaborative hypermedia information systems"

- Fundamental protocol behind "the web"
- Now part of most things we do on the Internet—so much more than web pages



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- Fundamental protocol behind "the web"
- Now part of most things we do on the Internet—so much more than web pages

But what is hypertext?

\dot{x}_A 70 languages \sim

Hypertext

Article Talk

From Wikipedia, the free encyclopedia

For the concept in semiotics, see Hypertext (semiotics).

Hypertext is text displayed on a computer display or other electronic devices with references (hyperlinks) to other text that the reader can immediately access.^[1] Hypertext documents are interconnected by hyperlinks, which are typically activated by a mouse click, keypress set, or screen touch. Apart from text, the term "hypertext" is also sometimes used to describe tables, images, and other presentational content formats with integrated hyperlinks. Hypertext is one of the key underlying concepts of the World Wide Web,^[2] where Web pages are often written in the Hypertext Markup Language (HTML). As implemented on the Web, hypertext enables the easy-to-use publication of information over the Internet.

Etymology [edit]

"(...)'Hypertext' is a recent coinage. 'Hyper-' is used in the mathematical sense of extension and generality (as in 'hyperspace,' 'hypercube') rather than the medical sense of 'excessive' ('hyperactivity'). There is no implication about <u>size</u>— a hypertext could contain only 500 words or so. 'Hyper-' refers to structure and not size."

- Theodor H. Nelson, Brief Words on the Hypertext 2, 23 January 1967

The English prefix "hyper-" comes from the Greek prefix "ὑπερ-" and means "over" or "beyond"; it has a common origin with the prefix "super-" which comes from Latin. It signifies the overcoming of the previous linear constraints of written text.



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Information mapping Topics and fields Business decision mapping • Data visualization

"As we may think", Vannevar Bush (1945)

"The human mind...operates by association. With one item in its grasp, it snaps instantly to the next ... in accordance with some intricate web of trails carried by the cells of the brain"

<u>"As we may think", Vannevar Bush (1945)</u>

"The human mind...operates by association. With one item in its grasp, it snaps instantly to the next ... in accordance with some intricate web of trails carried by the cells of the brain"

Defines the "Memex": "a device in which an individual stores all his books, records, and communications, and which is mechanized so that it may be consulted with exceeding speed and flexibility"



○ A https://en.wikipedia.org/wiki/HTTP

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Encrypted connections

Similar protocols

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From Wikipedia, the free encyclopedia			Ø	
(Redirected from Http)				
The Hypertext Transfer Protocol (HTTP) is an application layer protocol in the Internet protocol suite		НТТР		
communication for the World Wide Web, where hypertext documents include hyperlinks to other resources				

that the user can easily access, for example by a mouse click or by tapping the screen in a web browser. Development of HTTP was initiated by Tim Berners-Lee at CERN in 1989 and summarized in a simple document describing the behavior of a client and a server using the first HTTP version, named 0.9.^[2] That version was subsequently developed, eventually becoming the public 1.0.^[3]

Development of early HTTP Requests for Comments (RFCs) started a few years later in a coordinated effort by the Internet Engineering Task Force (IETF) and the World Wide Web Consortium (W3C), with work later moving to the IETF.

HTTP/1 was finalized and fully documented (as version 1.0) in 1996.^[4] It evolved (as version 1.1) in 1997 and then its specifications were updated in 1999, 2014, and 2022.^[5]

Its secure variant named HTTPS is used by more than 85% of websites.^[6] HTTP/2, published in 2015, provides a more efficient expression of HTTP's semantics "on the wire". As of April 2023, it is used by 39% of websites^[7] and supported by almost all web browsers (over 97% of users).^[8] It is also supported by

HTTP International RFC 1945 7 HTTP/1.0 standard RFC 9110 2 HTTP Semantics RFC 9111 Z HTTP Caching RFC 9112 Z HTTP/1.1 RFC 9113 2 HTTP/2 RFC 7541 7 HTTP/2: HPACK Header Compression RFC 8164 A HTTP/2: Opportunistic Security for HTTP: a protocol for distributing hypertext media (*and now so much more)

HTTP: a protocol for distributing hypertext media (*and now so much more)

Enables the World Wide Web (WWW): a distributed database of pages linked through HTTP

HTTP: a protocol for distributing hypertext media (*and now so much more)

Enables the World Wide Web (WWW): a distributed database of pages linked through HTTP

... now synonymous with with "The Internet" itself!

<u>Tim Berners-Lee</u>

- <u>1990</u>: First HTTP implementation – Tim Berners-Lee, CERN
- <u>1991</u>: HTTP/0.9: Fetching pages
- <u>1992</u>: HTTP/1.0: Client/server information, simple caching
- <u>1996</u>: HTTP/1.1
 - Extensive caching support
 - Host identification
 - Pipelined, persistent connections, ...



The first webserver

• <u>2015</u>: HTTP/2

– Main goal: reduce latency

• <u>2022</u>: HTTP/3

- Still: reduce latency
- Integrates security via TLS
- Replace transport layer with QUIC
- Already supported in >94% of browsers



http://httpwg.org/specs/rfc7540.html

How does "the web" work?

Webserver example.com

page.html <html> <title>hi</title> <h1>Welcome!</h1> </html>









Server returns response (in this case, with HTML)

Why so successful?

<u>Anyone</u> can host a website!

... just need a domain and a server

Why so successful?

<u>Anyone</u> can host a website!

... just need a domain and a server

Clients can easily find arbitrary pages, pages can easily link to others => content can grow very quickly

HTTP components

<u>Content</u>: objects (HTML, images, JSON, ...)

<u>Clients</u>: send <u>requests</u>, receive <u>response</u>

Servers: store content, or generate it

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<u>Servers</u>: store content, or generate it

Proxies/Middleboxes

- Placed between clients and servers
- Do extra stuff: caching, anonymization, logging, transcoding, filtering access

=> Important for scaling, modern browsing... more on this later

<u>How to find stuff?</u>

• So far: DNS: names for one or more hosts

– eg. cs.brown.edu

How do we ask for a specific *resource* from this host?

URL: Uniform Resource Locator

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- So far: DNS: names for one or more hosts
 - eg. cs.brown.edu

How do we ask for a specific *resource* from this host?

URL: Uniform Resource Locator

URLs: how we find stuff

https://cs.brown.edu/courses/csci1680/f23/policies/#late-policy

How to find stuff: URLs

protocol://[name@]hostname[:port]/directory/resource?k1=v1&k2=v2#tag

- Name: can identify a client
- *Hostname:* FQDN or IP address
- Port number: defaults to common protocol port (eg. 80, 22)
- Directory: path to the resource
- *Resource:* name of the object
- After that, various delimiters to specify further, common examples:
 - ?parameters are passed to the server for execution
 - #tag allows jumps to named tags within document

HTTP: the protocol

- Client-server protocol
- Protocol (but not data) in ASCII (before HTTP/2)
- Stateless
- Server typically listens on port 80 (or 443, with TLS)

• Server sends response, may close connection (client may ask it to say open)

Steps in HTTP^(1.0) Request

- Open TCP connection to server
- Send request
- Receive response
- TCP connection terminates
 - How many RTTs for a single request?
- You may also need to do a DNS lookup first!

> telnet www.cs.brown.edu 80 Trying 128.148.32.110... Connected to www.cs.brown.edu. Escape character is '^]'. GET / HTTP/1.0 HTTP/1.1 200 OK Date: Thu, 24 Mar 2011 12:58:46 GMT Server: Apache/2.2.9 (Debian) mod_ssl/2 Last-Modified: Thu, 24 Mar 2011 12:25:2

Server: Apache/2.2.9 (Debian) mod_ssl/2.2.9 OpenSSL/0.9.8g
Last-Modified: Thu, 24 Mar 2011 12:25:27 GMT
ETag: "840a88b-236c-49f3992853bc0"
Accept-Ranges: bytes
Content-Length: 9068
Vary: Accept-Encoding
Connection: close
Content-Type: text/html

```
<!DOCTYPE html PUBLIC "-//W3C//DTD XHTML 1.0 Strict//EN"
    "http://www.w3.org/TR/xhtml1/DTD/xhtml1-strict.dtd">
```

HTTP Request

<u>Method</u>:

- GET: current value of resource, run program
- POST: update a resource, provide input for a program. . .

<u>Headers</u>: useful info about request

- E.g., desired language, text encoding



Sample Browser Request

GET / HTTP/1.1 Host: localhost:8000 User-Agent: Mozilla/5.0 (Macinto ... Accept: text/xml,application/xm ... Accept-Language: en-us,en;q=0.5 Accept-Encoding: gzip,deflate Accept-Charset: ISO-8859-1,utf-8;q=0.7,*;q=0.7 (empty line)

In your browser: right click => Inspect element => Network

HTTP Responses

Status codes to indicate something about the result

- 1xx: Information e.g, 100 Continue
- 2xx: Success e.g., 200 OK
- 3xx: Redirection e.g., 302 Found (elsewhere),
- 4xx: Client Error e.g., 403 Forbidden, 404 Not Found
- 5xx: Server Error e.g, 503 Service Unavailable

status	version	statu	s code	phrase ←		
	header fie	ld name	value	←↓		
headers					ļ	
	header fie	ld name	value	←↓		
blank line	←↓					
body						,, ,,

HTTP is Stateless

- Each request/response treated independently
- Servers not required to maintain state

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But...

- Most applications need persistent state
- E.g., shopping cart, web-mail, usage tracking, (most sites today!)

HTTP Cookies

- Client-side state maintenance
 - Client stores small state on behalf of server
 - Sends request in future requests to the server
 - Cookie value is meaningful to the server (e.g., session id)
- Can provide authentication



Anatomy of a Web Page

- HTML content
- A number of additional resources
 - Images
 - Scripts
 - Frames
- Browser makes one HTTP request for each object
 - Course web page: 14 objects
 - Modern web pages: hundreds of objects



of **Computer Science**

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Welcome to the Brown University Computer Science Department Web. Information here is organized into broad categories, which are summarized in the icon bar, above. If you are visiting for th or exploring, the rest of this page offers some details about what you'll find.

If you are visiting us in person, you'll need directions to the CIT building. If not, perhaps you just need our address, phone, fax or other vital statistics.

Calendar of Events

Talks, conferences and soirces both at Brown and elsewhere are described.

Programs of Study

Undergraduate concentration requirements and the masters and phd programs are described, accompanied by the relevant forms, brochures and pointers to related information elsewhere.

Research Groups

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Active research areas in computer science at Brown include graphics, geometric computing, object-oriented databases, artificial intelligence and robotics. Each group maintains a home pag describing their research and activities and links to relevant publications.

Publications

The Department publishes brochures, technical reports, a newsletter, *conduit!*, and, for locals, house rules.

Courses

Many courses taught using the Department's facilities have home pages, which provide information useful to students taking them.



Modern web pages and HTTP

- Web APIs: HTTP response/requests are a standard way to ask for anything
- Modern web pages: use Javascript to make lots of requests without reloading page
 - And can use APIs for all kinds of other stuff

Example: Github public API

```
$ curl https://api.github.com/users/ndemarinis
  "login": "ndemarinis",
  "id": 1191319,
  "node_id": "MDQ6VXNlcjExOTEzMTk=",
  "avatar_url": "https://avatars.githubusercontent.com/u/1191319?v=4",
  "gravatar id": "",
  "url": "https://api.github.com/users/ndemarinis",
  "type": "User",
  "site admin": false,
  "name": "Nick DeMarinis",
  "blog": "https://vty.sh",
  "twitter_username": null,
  "public repos": 10,
```

HTTP

> telnet www.cs.brown.edu 80 Trying 128.148.32.110... Connected to www.cs.brown.edu. Escape character is '^]'. GET / HTTP/1.0 HTTP/1.1 200 OK Date: Thu, 24 Mar 2011 12:58:46 GMT Server: Apache/2.2.9 (Debian) mod ssl/2.2.9 OpenSSL/0.9.8g Last-Modified: Thu, 24 Mar 2011 12:25:27 GMT ETag: "840a88b-236c-49f3992853bc0" Accept-Ranges: bytes Content-Length: 9068 Vary: Accept-Encoding Connection: close Content-Type: text/html

```
<!DOCTYPE html PUBLIC "-//W3C//DTD XHTML 1.0 Strict//EN"
    "http://www.w3.org/TR/xhtml1/DTD/xhtml1-strict.dtd">
```

HTTP: What matters for performance?

Depends on type of request

- Lots of small requests (objects in a page)
- Some big requests (large download or video)

Small Requests

- Latency matters
- RTT dominates
- Major steps:
 - DNS lookup (if not cached)
 - Opening a TCP connection
 - Setting up TLS (optional, but now common)
 - Actually sending the request and receiving response

How can we reduce the number of connection setups?

- Keep the connection open and request all objects serially
 - Works for all objects coming from the same server
 - Which also means you don't have to "open" the window each time

Persistent connections (HTTP/1.1)

Browser Request

GET / HTTP/1.1 Host: localhost:8000 User-Agent: Mozilla/5.0 (Macinto ... Accept: text/xml,application/xm ... Accept-Language: en-us,en;q=0.5 Accept-Encoding: gzip,deflate Accept-Charset: ISO-8859-1,utf-8;q=0.7,*;q=0.7 Keep-Alive: 300 Connection: keep-alive

Small Requests (cont)

- Second problem is that requests are serialized
 - Similar to stop-and-wait protocols!
- Two solutions
 - Pipelined requests (similar to sliding windows)
 - Parallel Connections
 - Browsers implement this differently—see "Inspect element"
 - How are these two approaches different?

HTTP/2

- Adds more options to trade off:
- Multiplexed streams on same connection
 - Plus stream weights, dependencies
- No head of line blocking!
 - But what happens if there is packet loss?

https://www.twilio.com/blog/2017/10/http2-issues.html

HTTP/3

- Mapping of HTTP semantics onto QUIC
 - E.g., QUIC already implements multiple streams, and HTTP doesn't need to do it
- QUIC: Another transport-layer protocol, intended to replace TCP
 - RFC9000
 - Same goals as TCP, but...
 - Integrates security by default (TLS, next class)
 - Supports multiple streams at once
 - Various tricks to reduce message size and latency
- By moving multiplexing into the transport layer, can do so in a way that benefits HTTP (no head of line blocking!)

Comparison: QUIC's handshake



Larger Objects

- Problem is throughput in bottleneck link
- Solution: HTTP Proxy Caching
 - Also improves latency, and reduces server load



How to Control Caching?

- Server sets options
 - Expires header
 - No-Cache header
- Client can do a conditional request:
 - Header option: if-modified-since
 - Server can reply with 304 NOT MODIFIED

Caching

- Where to cache content?
 - Client (browser): avoid extra network transfers
 - Server: reduce load on the server
 - Service Provider: reduce external traffic



Caching

- Why caching works?
 - Locality of reference:
 - Users tend to request the same object in succession
 - Some objects are popular: requested by many users



How well does caching work?

- Very well, up to a point
 - Large overlap in requested objects
 - Objects with one access place upper bound on hit ratio
 - Dynamic objects not cacheable*
- Example: Wikipedia
 - About 400 servers, 100 are HTTP Caches (Squid)
 - 85% Hit ratio for text, 98% for media

* But can cache portions and run special code on edges to reconstruct

Reverse Proxies

Close to the server

– Also called Accelerators



Forward Proxies

Typically done by ISPs or Enterprises

- Reduce network traffic and decrease latency
- May be transparent or configured



Content Distribution Networks

- Integrate forward and reverse caching
 - One network generally administered by one entity
 - E.g. Akamai
- Provide document caching
 - Pull: result from client requests
 - Push: expectation of high access rates to some objects
- Can also do some processing
 - Deploy code to handle some dynamic requests
 - Can do other things, such as transcoding

Example CDN



How Akamai works

Akamai has cache servers deployed close to clients

- Co-located with many ISPs
- Challenge: make same domain name resolve to a proxy close to the client
- Lots of DNS tricks. BestBuy is a customer
 - Delegate name resolution to Akamai (via a CNAME)

DNS Resolution

dig www.bestbuy.com						
;; ANSWER SECTION:						
www.bestbuy.com. 3600	IN	CNAME	www.bestbuy.com.edgesuite.net.			
www.bestbuy.com.edgesuite.net. 21600 IN			CNAME	all05.b.akamai.net.		
a1105.b.akamai.net.	20	IN	A	198.7.236.235		
a1105.b.akamai.net.	20	IN	A	198.7.236.240		
;; AUTHORITY SECTION:						
b.akamai.net.	1101	IN	NS	nlb.akamai.net.		
b.akamai.net.	1101	IN	NS	nOb.akamai.net.		
;; ADDITIONAL SECTION:						
nOb.akamai.net.	1267	IN	A	24.143.194.45		
nlb.akamai.net.	2196	IN	A	198.7.236.236		

- n1b.akamai.net finds an edge server close to the client's local resolver
 - Uses knowledge of network: BGP feeds, traceroutes. *Their secret sauce...*

Example

From Brown

dig www.bestbuy.com
;; ANSWER SECTION:
www.bestbuy.com. 3600 IN CNAME www.bestbuy.com.edgesuite.net.
www.bestbuy.com.edgesuite.net. 21600 IN CNAME al105.b.akamai.net.
al105.b.akamai.net. 20 IN A 198.7.236.235
al105.b.akamai.net. 20 IN A 198.7.236.240
- Ping time: 2.53ms

	<u>From Berkeley, CA</u>				
a1105.b.akamai.net.	20	IN	А	198.189.255.200	
all05.b.akamai.net.	20	IN	А	198.189.255.207	
 Ping time: 3.20ms 					

dig www.bestbuy.com
;; QUESTION SECTION:
;www.bestbuy.com. IN A

;; ANSWER SECTION: www.bestbuy.com. 2530 IN CNAME www.bestbuy.com.edgekey.net. www.bestbuy.com.edgekey.net. 85 IN CNAME e1382.x.akamaiedge.net. e1382.x.akamaiedge.net. 16 IN A 104.88.86.223

;; Query time: 6 msec ;; SERVER: 192.168.1.1#53(192.168.1.1) ;; WHEN: Thu Nov 16 09:43:11 2017 ;; MSG SIZE rcvd: 123

traceroute to 104.88.86.223 (104.88.86.223), 64 hops max, 52 byte packets

- 1 router (192.168.1.1) 2.461 ms 1.647 ms 1.178 ms
- 2 138.16.160.253 (138.16.160.253) 1.854 ms 1.509 ms 1.462 ms
- 3 10.1.18.5 (10.1.18.5) 1.886 ms 1.705 ms 1.707 ms
- 4 10.1.80.5 (10.1.80.5) 4.276 ms 6.444 ms 2.307 ms
- 5 lsb-inet-r-230.net.brown.edu (128.148.230.6) 1.804 ms 1.870 ms 1.727 ms
- 6 131.109.200.1 (131.109.200.1) 2.841 ms 2.587 ms 2.530 ms
- 7 host-198-7-224-105.oshean.org (198.7.224.105) 4.421 ms 4.523 ms 4.496 ms
- 8 5-1-4.bear1.boston1.level3.net (4.53.54.21) 4.099 ms 3.974 ms 4.290 ms
- 9 * ae-4.r00.bstnma07.us.bb.gin.ntt.net (129.250.66.93) 4.689 ms 4.109 ms
- 10 ae-6.r24.nycmny01.us.bb.gin.ntt.net (129.250.4.114) 8.863 ms 10.205 ms 10.477 ms
- 11 ae-1.r08.nycmny01.us.bb.gin.ntt.net (129.250.5.62) 9.298 ms ae-1.r07.nycmny01.us.bb.gin.ntt.net (129.250.3.181) 10.008 ms 8.677 ms
- 12 ae-0.a00.nycmny01.us.bb.gin.ntt.net (129.250.3.94) 8.543 ms 7.935 ms ae-1.a00.nycmny01.us.bb.gin.ntt.net (129.250.6.55) 9.836 ms
- 13 a104-88-86-223.deploy.static.akamaitechnologies.com (104.88.86.223) 9.470 ms 8.483 ms 8.738 ms

dig www.bestbuy.com @109.69.8.51 e1382.x.akamaiedge.net. 12 IN A 23.60.221.144 traceroute to 23.60.221.144 (23.60.221.144), 64 hops max, 52 byte packets 1 router (192.168.1.1) 44.072 ms 1.572 ms 1.154 ms 2 138.16.160.253 (138.16.160.253) 2.460 ms 1.736 ms 2.722 ms 3 10.1.18.5 (10.1.18.5) 1.841 ms 1.649 ms 3.348 ms 4 10.1.80.5 (10.1.80.5) 2.304 ms 15.208 ms 2.895 ms 5 lsb-inet-r-230.net.brown.edu (128.148.230.6) 1.784 ms 4.744 ms 1.566 ms 6 131.109.200.1 (131.109.200.1) 3.581 ms 5.866 ms 3.238 ms 7 host-198-7-224-105.oshean.org (198.7.224.105) 4.288 ms 6.218 ms 8.332 ms 8 5-1-4.bear1.boston1.level3.net (4.53.54.21) 4.209 ms 6.103 ms 5.031 ms 9 ae-4.r00.bstnma07.us.bb.gin.ntt.net (129.250.66.93) 3.982 ms 5.824 ms 4.514 ms 10 ae-6.r24.nycmny01.us.bb.gin.ntt.net (129.250.4.114) 9.735 ms 12.442 ms 8.689 ms 11 ae-9.r24.londen12.uk.bb.gin.ntt.net (129.250.2.19) 81.098 ms 81.343 ms 81.120 ms 12 ae-6.r01.mdrdsp03.es.bb.gin.ntt.net (129.250.4.138) 102.009 ms 110.595 ms 103.010 ms 13 81.19.109.166 (81.19.109.166) 99.426 ms 93.236 ms 101.168 ms 14 a23-60-221-144.deploy.static.akamaitechnologies.com (23.60.221.144) 94.884 ms 92.77 ms 93.281 ms

Other CDNs

- Akamai, Limelight, Cloudflare
- Amazon, Facebook, Google, Microsoft
- Netflix
- Where to place content?
- Which content to place? Pre-fetch or cache?