CSCI 1680 Physical Layer, Link Layer I

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Based partly on lecture notes by Rodrigo Fonseca, David Mazières, Phil Levis, John Jannotti

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In TCP, conn.Read returns when ANY AMOUNT of data is available => MAY NOT be the amount you want

Solution: keep reading until you get the amount of data that you want io.ReadFull

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SENDING DATA W/ TCP

Image: A matrix in the image

Image: Conv.wHfE()

Image: Conv.white

Image: Conv.w

In Go: io.ReadFull

TCP is designed to provide a STREAM of ordered data => it doesn't care about the separation of individual messages

WHAT HAPPENS IN YOU DON'T KNOW THE SIZE OF THE MESSAGE? => PROTOCOL NETERS TO BE SET UP SO THAT YOU CAN ALWAYS FIGURE BUT NOW MUCH DATA TO READ NEXT. TYPE LENGTH STRING F6,

- MESSLOE TIPE => SIZE

> LENGTH OF DATH TO FOLLOW



• Two more things on sockets

Physical/Link layer: how to connect two things
 => Inherent properties of *real* networks

Layers, Services, Protocols

	Application	Service: user-facing application. Application-defined messages
	Transport	Service: multiplexing applications Reliable byte stream to other node (TCP), Unreliable datagram (UDP)
	Network	Service: move packets to any other node in the network IP: Unreliable, best-effort service model
12	Link	Service: move frames to other node across link. May add reliability, medium access control
_)	Physical	Service: move bits to other node across link

Physical Layer (Layer 1)

Specifies three things:

- Physical medium; WIPI, ETHERNET, CELLULAR...
 Signaling/modulation: How DO I SEND
 Encoding:
- Encoding \











Physical Layer (Layer 1)

Specifies three things:

- Physical medium: cable, fiber, wireless frequency
- Signaling/modulation: how to transmit/receive
- Encoding: how to get meaningful data









Why should we care?

This is the line between electrical engineering and computer science

Helpful to understand challenges involved => How design/limitations affect our systems

Also: Learn important principles we'll use elsewhere



Send/receive data over some kind of medium Sender encodes message in some format, sends it "over the wire" Receiver decodes message (knowing the format) => recover the message

Why is this hard?

- Sharing channel: interference from other devices
- Noise
- Physical distance (attenuation)
- Energy usage
- Security
- •

=> Every medium has its own characteristics, and problems

We don't need to know the details.

However, there are some key takeways to help understand the challenges and implications



1. All media have fixed <u>bandwidth</u> => fixed "space" to transmit information

2. Sending data takes time! => <u>latency</u>

3. All media have (some) <u>errors</u> => how to deal with them?

Bandwidth

Bandwidth: set of frequencies that a channel (medium) can propagate wall $(\mathcal{A}_{\mathcal{V}})$

PHONE/DIALOP: 8kHz

 Creates a fixed "space" in which data can be transmitted =>Wires: defined by physical properties ⇒ Wireless: frequency ranges are regulated

Bandwidth gives an upper bound on <u>throughput</u> => amount of data we can send per time (bits / second)

Bandwidth: frequencies that a channel propagates well (Most signals made up of different frequencies)

Creates a fixed "space" in which data can be transmitted
 =>Wires: defined by physical properties
 ⇒ Wireless: frequency ranges are regulated

Upper bound on *throughput*: amount of data we can send per time (bits per second)

UNITED

STATES FREQUENCY ALLOCATIONS

THE RADIO SPECTRUM



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Early IEEE 802.11 (Wifi) channel bandwidth



2,16Hz: "UNIL CTUSTO BIND"

Early IEEE 802.11 (Wifi) channel bandwidth



Wi-Fi generations				
Generation	IEEE standard	Adopted	Maximum link rate (Mbit/s)	Radio frequency (GHz)
Wi-Fi 7	802.11be	(2024)	1376 to 46120	2.4/5/6
Wi-Fi 6E	902 11 ox	2020	E74 to 0609[41]	6 ^[42]
Wi-Fi 6	002.11ax	2019	574 10 9000	2.4/5
Wi-Fi 5	802.11ac	2014	433 to 6933	5 ^[43]
Wi-Fi 4	802.11n	2008	72 to 600	2.4/5
(Wi-Fi 3)*	802.11g	2003	6 to 54	2.4
	802.11a	1999	010 54	5
(Wi-Fi 2)*	802.11b	1999	1 to 11	2.4
(Wi-Fi 1)*	802.11	1997	1 to 2	2.4
*(Wi-Fi 1, 2, and 3 are by retroactive inference) [44][45][46][47][48]				

IEEE 802,11

How to actually send data?

(Within a limited bandwidth)

How to actually send stuff?



One way: Use Carriers

Start with a carrier frequency, modulate it to encode data:









This can get more complex...

Lots of engineering you can do

- Multiple carriers/frequencies
- Adjust amplitude, phase

• Clever ways to avoid errors



A good animation on Wikipedia

QAM: Quadrature Amplitude Modulation

Example: Quadrature Amplitude Modulation (QAM)

256-QAM Constellation



Modulation schemes in action

• <u>https://www.youtube.com/watch?v=vvr9AMWEU-c</u>

Sounds great, right?

• Problem: noise limits the number of modulation levels (M)

MODULATION MOLEVELS 0) 00

Sounds great, right?

• Problem: noise limits the number of modulation levels (M)

Shannon's Law: $C = B \log_2(1 + S/N)$

- C: channel capacity (throughput) in bits/second
- B : bandwidth in Hz
- S, N: average signal, noise power

The amount of data we can fit in a channel is limited by the bandwidth, and the amount of noise in the medium

Sounds great, right?

• Problem: noise limits the number of modulation levels (M)

Shannon's Law: $C = B \log_2(1 + S/N)$

- C: channel capacity (throughput) in bits/second
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- S, N: average signal, noise power

Takeaway: fundamental limit on how much data we can fit into a fixed channel, based on noise => For any medium, designers create encodings to try and maximize throughput

Medium	Bandwidth	Throughput	
Dialup	8 kHz	56 Kbit/s	9
Early Wifi (802.11g)	20 MHz	54 Mbit/s	PANGÉ
Modern Wifi (802.11ax)	20-40 MHz	Up to 9 Gbps	INTERFERENCE
Ethernet	62.5 MHz_ (1Gbps version)	1Gbit/s (common) Up to 100Gbps	10-12
3G cellular	Depends on carrier	2 Mbit/s	0
5G cellular	Depends on carrier	> 1 GBps	

=> Does this mean wifi is the best?

Latency

Sending data takes time!

- Latency: time between sending data and when data arrives (somewhere)
 Sending data is not instantaneous (for many reasons)
- Multiple components => many definitions, depending on what we're measuring A B $\pm \int \int \int \int \frac{1}{\sqrt{2}} dx$ $-\frac{1}{\sqrt{2}} \int \int \frac{1}{\sqrt{2}} dx$

How to think about latency



Processing delay: Time for data to be sent by a node
 => Time for OS to decide to send data

Transmission delay: Time for transceiver to actually send bits on wire

Propagation delay: Time for data to propagate across wire => speed of light

Queuing delay: time data spends in buffers waiting to be sent (due to congestion, etc.)

How to think about latency



How to think about latency

- <u>Processing delay</u> at the node: per message computation
- <u>Queuing delay</u>: time spent waiting in buffers
- Transmission delay: sending out the actual data
 Size/Bandwidth
- <u>Propagation delay</u>: time for bits to actually go out on the wire
 - Upper bound?
 - Depends on media, ultimate upper bound is speed of light

Round trip time (RTT): time between request and response

When we design protocols, can think about performance based on number of RTTs



Next few pages are extra material we'll discuss later!

Sending Frames Across



Throughput: bits / s

Which matters most, bandwidth or delay?

- How much data can we send during one RTT?
- E.g., send request, receive file



Often: For small transfers, latency more important, for bulk, throughput more important

Performance Metrics

- <u>Throughput</u>: Number of bits received/unit of time
 e.g. 100 Mbps
- <u>Goodput</u>: Useful bits received per unit of time
- Latency: How long for message to cross network
- <u>Jitter</u>: Variation in latency

Dealing with errors

Error Detection

• Basic idea: use a checksum

- Compute small check value, like a hash of packet

Error Detection

• Basic idea: use a checksum

- Compute small check value, like a hash of packet

- Good checksum algorithms
 - Want several properties, e.g., detect any single-bit error
 - Details later

=> Not all protocols do this. Why?

Approximation of a Square Wave



Graphs from Dr. David Alciatore, Colorado State University